Higher Groups in Homotopy Type Theory

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September 24, 2019

Part 1: Homotopy Type Theory

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- To identify Hⁿ(Sⁿ; Z) with Z, we must choose an orientation for the n-sphere Sⁿ.
- To identify the natural number n such that $\pi_4(S^3)$ is isomorphic to $\mathbb{Z}_{/n}$ with the number 2, we need to prove that n equals 2.

Type Theory

A type is a type of mathematical thing.

Type theory gives *rules* for making new types and new terms of them. It is a full foundation of mathematics, from scratch.

> a term : its type a : A 3 : \mathbb{N} \mathbb{N} : Set T_pM : Vect_R Vect_R : Type $K(\mathbb{Z}; 2)$: Type

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- For n and m : \mathbb{N} , then n = m has an element if and only if n equals m
 - there is at most one way to identify two natural numbers.

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Axiom (Univalence)

If X and Y are types, then X = Y is the type of equivalences of X with Y.

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Types can depend on elements of others types.

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• A vector field is $v : (p : M) \to T_pM$.

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 $(\mathsf{a}:\mathsf{A})\times\mathsf{B}(\mathsf{a})$

is the type of pairs (a, b) with a : A and b : B(a).

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Definition

For a function $f : A \rightarrow B$ and b : B, its *fiber* is the type

 $fib_f(b) :\equiv (a : A) \times (f(a) = b)$

together with the map $(a, p) \mapsto a : fib_f(b) \to A$.

Contractible Types and the Substitution Lemma

Definition

A center of contraction for a type A is an element c : A such that for every other element a : A, we have an identification of a with c.

$$\mathbf{Contr}(\mathsf{A}) :\equiv (\mathsf{c} : \mathsf{A}) \times ((\mathsf{a} : \mathsf{A}) \to (\mathsf{a} = \mathsf{c}))$$

Definition

A map $f : A \rightarrow B$ is an *equivalence* if its fibers are contractible:

 $\textbf{Equiv}(f):\equiv (y:Y) \rightarrow \textbf{Contr}(fib_f(y))$

Contractible Types and the Substitution Lemma

Lemma (Singleton Lemma, UFP)

For any type A and element c : A, the singleton type ("based path type")

 $(a:A) \times (a=c)$

is contractible.

Lemma (Substitution Lemma, UFP) Suppose $B : A \rightarrow Type$ and c : Contr(A) is a center of contraction for A. Then

 $(a:A) \times B(a) = B(c).$

Corollary (Substitution Lemma)

Let c : A and suppose that $B : A \rightarrow Type$. Then

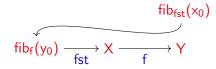
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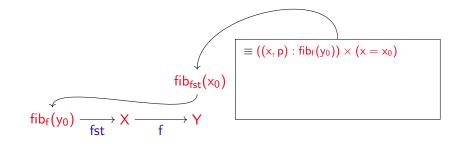
$$\mathsf{fib}_{\mathsf{f}}(\mathsf{y}_0) \xrightarrow{} \mathsf{fst} \mathsf{X} \xrightarrow{} \mathsf{f} \mathsf{Y}$$

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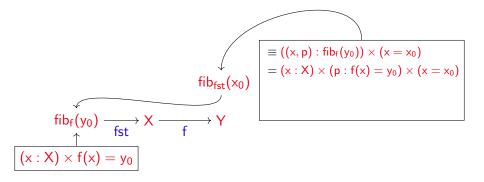
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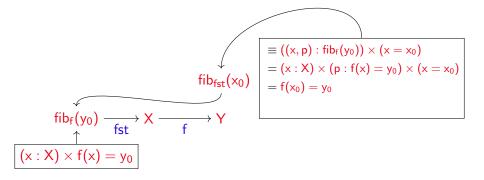
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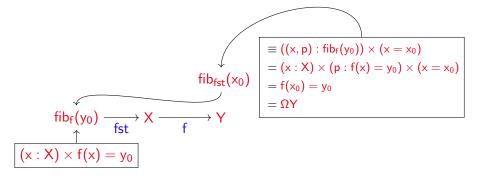
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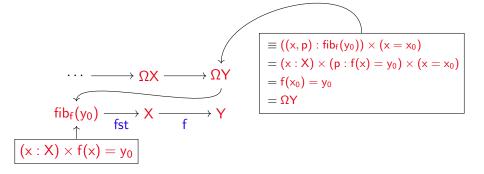
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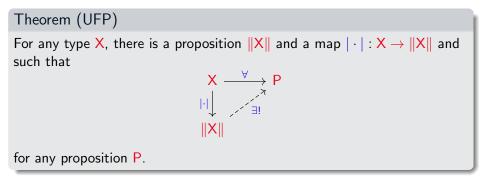
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- A groupoid is a type G such that for any a, b : G, a = b is a set.
- . . .
- An n-type is a type X such that for any a, b : X, a = b is an (n 1)-type (with -2-types being contractible).

Truncation

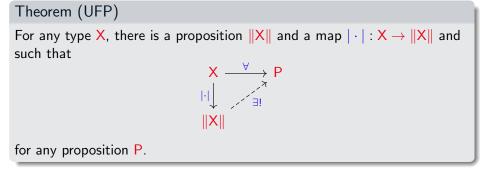
Theorem (UFP) For any type X, there is a proposition ||X|| and a map $|\cdot| : X \rightarrow ||X||$ and such that $\begin{array}{c} X \xrightarrow{\forall} P \\ |\cdot| \downarrow & \swarrow^{\pi} \\ ||X|| \end{array}$ for any proposition P.

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- If c : A, then (a : A) × ||a = c|| is all elements of A which are identifiable with c.

Part 2: Higher Groups

A group is the collection of symmetries of an object.

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• The dihedral group D_{2n} is the type of symmetries of an n-gon in the plane.

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 Let BGL_n(ℝ) :≡ (V : Vect_ℝ) × ||V = ℝⁿ|| be the type of n-dimensional vector spaces, pointed at ℝⁿ. Then

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• Let $BGL_n(\mathbb{R}) :\equiv (V : Vect_{\mathbb{R}}) \times ||V = \mathbb{R}^n||$ be the type of n-dimensional vector spaces, pointed at \mathbb{R}^n . Then

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Definition

For a : A, define

 $\mathsf{BAut}_\mathsf{A}(\mathsf{a}) :\equiv (\mathsf{b} : \mathsf{A}) \times \|\mathsf{b} = \mathsf{a}\|$

be the type of things *identifiable* with a, pointed at a. Then

 $Aut_A(a) \equiv a =_A a = \Omega BAut_A(a)$

$\infty ext{-Groups}$

Definition

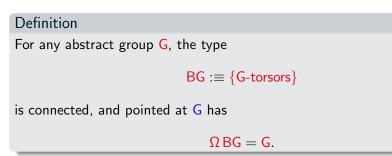
An ∞ -group is a type G identified with the loops of a pointed, connected type BG (the *delooping* of G).

 $\begin{aligned} & \infty \text{Group} :\equiv (\mathsf{G} : \mathbf{Type}) \times (\mathsf{BG} : \mathbf{Type}_*^{>0}) \times \mathsf{G} = \Omega \, \mathsf{BG} \\ & = \mathbf{Type}_*^{>0}. \end{aligned}$

Definition

An abstract group is a set ${\sf G}$ equipped with associative, unital, and invertible binary operation.

Groups are ∞ -Groups



A homomorphism of ∞ -groups is a pointed map between their deloopings:

 $\mathsf{BG} \, \cdot \to \, \mathsf{BH} :\equiv (\mathsf{f} : \mathsf{BG} \to \mathsf{BH}) \times (\mathsf{f}(\mathsf{pt}_{\mathsf{BG}}) = \mathsf{pt}_{\mathsf{BH}})$

- Consider $f \equiv X \mapsto X + 1$: BAut(n) $\cdot \rightarrow$ BAut(n + 1),
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- Consider $f \equiv V \mapsto \Lambda^n V : BGL_n(\mathbb{R}) \to BGL_1(\mathbb{R})$,
 - pointed by $e_1 \wedge \cdots \wedge e_n \mapsto 1 : \Lambda^n \mathbb{R}^n = \mathbb{R}$.

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 - Then $\Omega f = det$ takes the determinant.

Schreier Theory: Classifying Extensions

Lemma

For groups F and G the type of extensions $F \rightarrow E \rightarrow G$ of G by F is equivalent to the type of fiber sequences $BF \rightarrow BE \rightarrow BG$.

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Theorem (Schreier Theory for ∞ -groups)

Let F and G be ∞ -groups, then

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Let F and G be ∞ -groups, then

 $\{Extensions of G by F\} = BG \cdot \rightarrow BAut(BF)$

Corollary

For any group G,

 $\{\text{Extensions of } \mathbb{Z} \text{ by } G\}_{/\cong} = \textbf{Out}(G)$

References

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